

Fighting Fraud with Graph Databases

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Smals Research

Devoxx – Nov 8th, 2017

Outline

Motivations

Graph databases

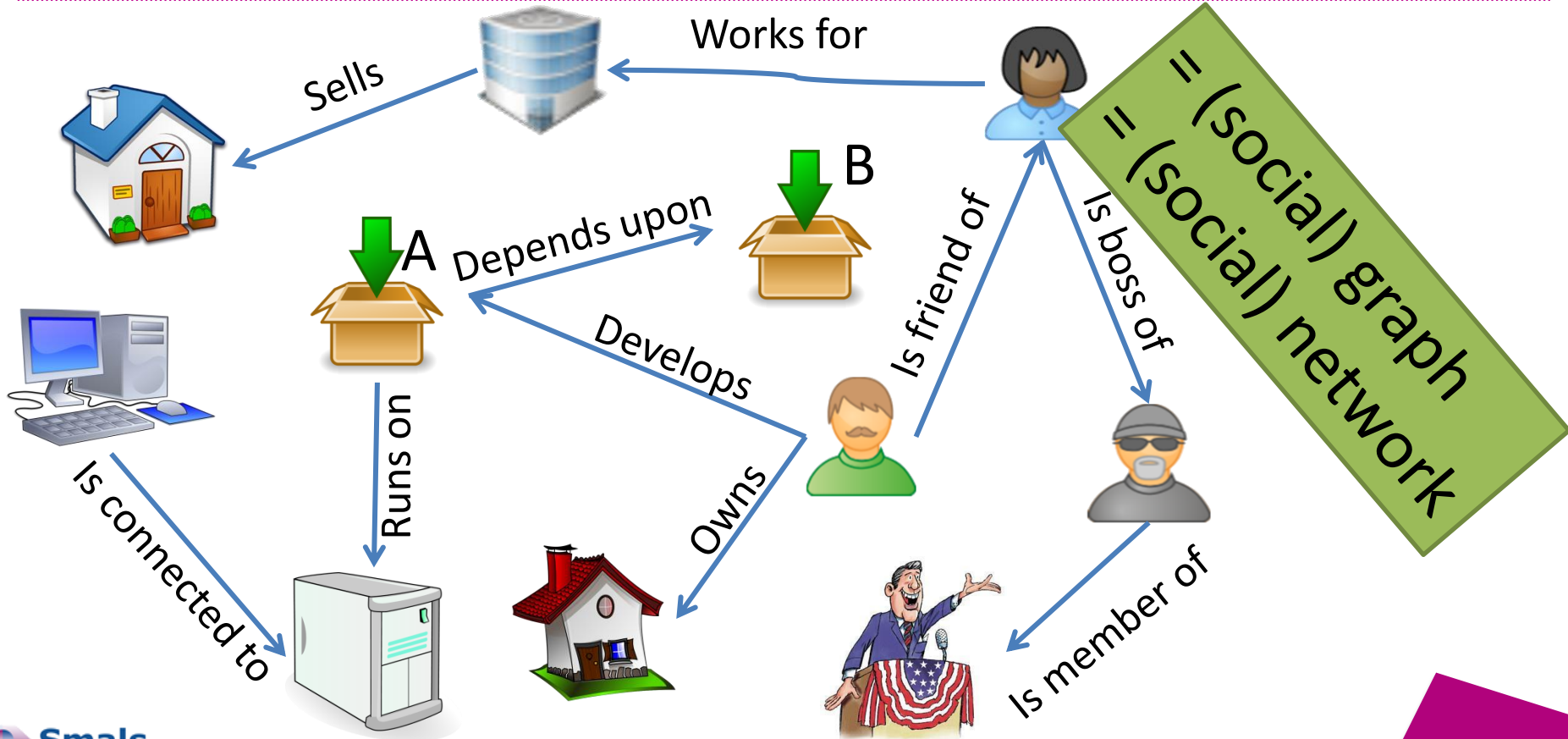
Neo4j

Use case : Social dumping

Conclusions

MOTIVATIONS

Motivations



Starting points

- Graphs excel in model fraudulent/criminal behavior
- RDMS perform poorly in representing graphs (i.e., relationships between entities)
- Graph databases propose a new database model, allowing to query graphs
- RBMS focus on entities, Graph DB focus on relationships

GRAPH DATABASES

RDBMS and relationships

Workers
ID
Name
Employer_ID

Companies
ID
Name

Smals' Employees?

How relationships are implemented

What concerns developers

```
SELECT Workers.Name
```

```
FROM Workers
```

```
JOIN Companies
```

```
ON Workers.Employer ID = Companies.ID
```

```
WHERE Companies.Name = 'Smals'
```

RDBMS : Relational DataBase Management System

RDBMS and relationships

Workers
ID
Name

Works_for
Worker_ID
Company_ID
From_date

Companies
ID
Name

Smals' Employees?

How relationships are implemented

What concerns developers

```
SELECT Workers.Name
FROM Workers
JOIN Works_for
ON Workers.ID = Works_for.Worker_ID
JOIN Companies
ON Works_for.Company_ID = Companies.ID
WHERE Companies.Name = 'Smals'
```


RDBMS and relationships

ID	Name
1	Alice
2	Bob
3	Camille
4	Zoé

Liker_ID	Liked_ID
1	3
2	1
2	3
3	2
3	4

```
SELECT p1.Name
```

```
FROM People p1
```

```
JOIN Likes
```

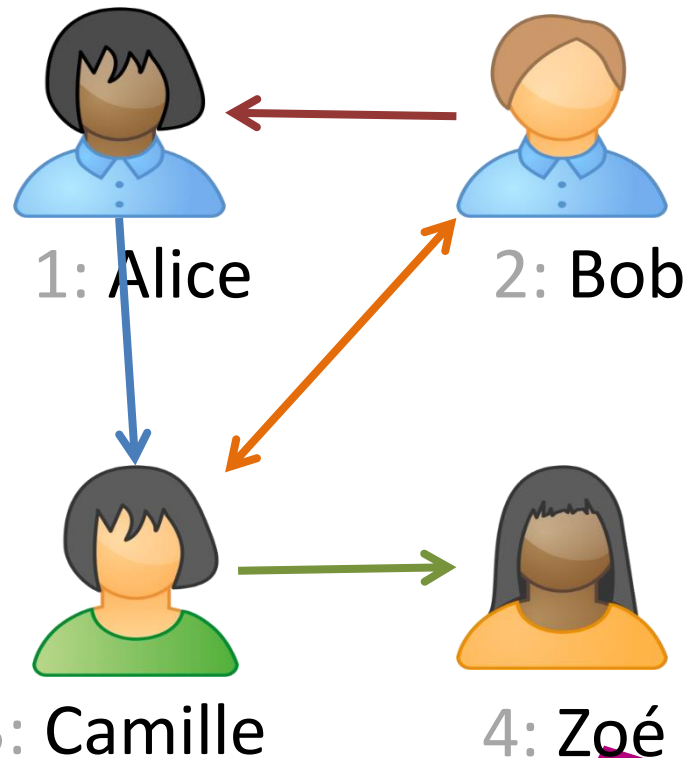
```
ON Likes.Liked_ID = p1.ID
```

```
JOIN People p2
```

```
ON Likes.Liker_ID = p2.ID
```

```
WHERE p2.Name = "Bob"
```

Who Bob likes?



RDBMS and relationships

ID	Name
1	Alice
2	Bob
3	Camille
4	Zoé

Liker_ID	Liked_ID
1	3
2	1
2	3
3	2
3	4

```
SELECT p1.Name
```

```
FROM People p1
```

```
JOIN Likes l1
```

```
ON l1.Liked_ID = p1.ID
```

```
JOIN Likes l2
```

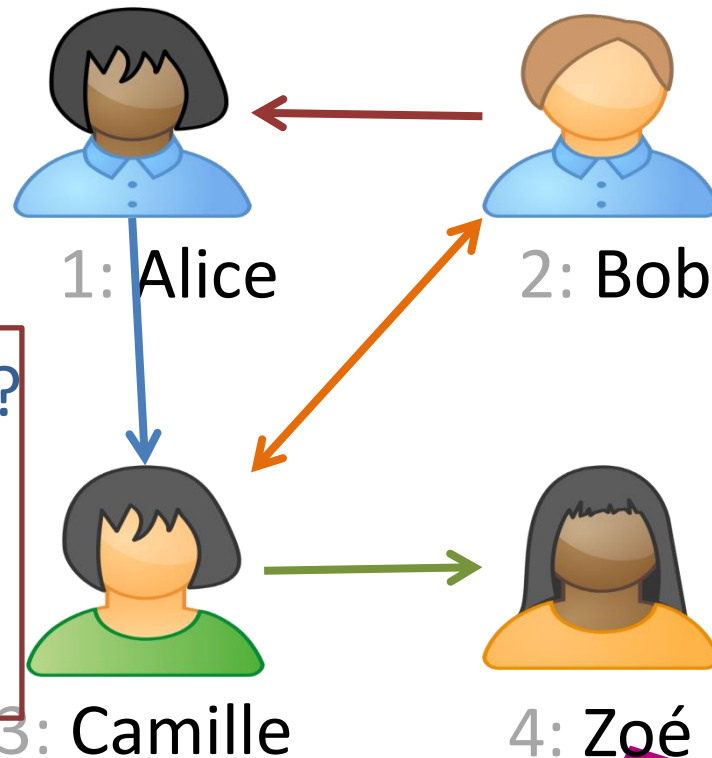
```
ON l1.Liker_ID = l2.Liked_ID
```

```
JOIN People p2
```

```
ON l2.Liker_ID = p2.ID
```

```
WHERE p2.Name = "Bob"
```

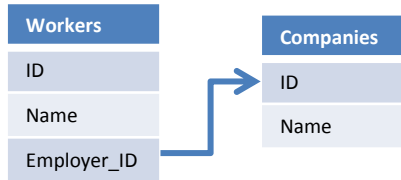
Who Bob (likes)² ?



RDBMS and relationships: limitations

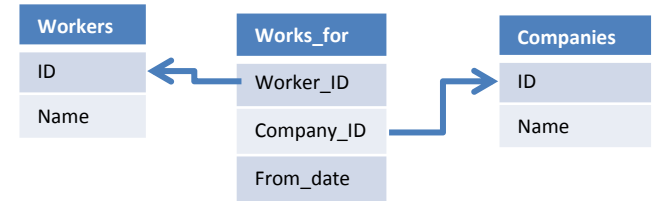
1-n relationships

Foreign key:



m-n relationships

Join table:



- Transforms the role of **attributes**
- No type « Key »
- Integrity constraints possible (extern to tables)
- Relationship structure to be given to each request
- Join : Heavy in writing code and in execution time
- Transforms the role of **tables**
- No distinguishable

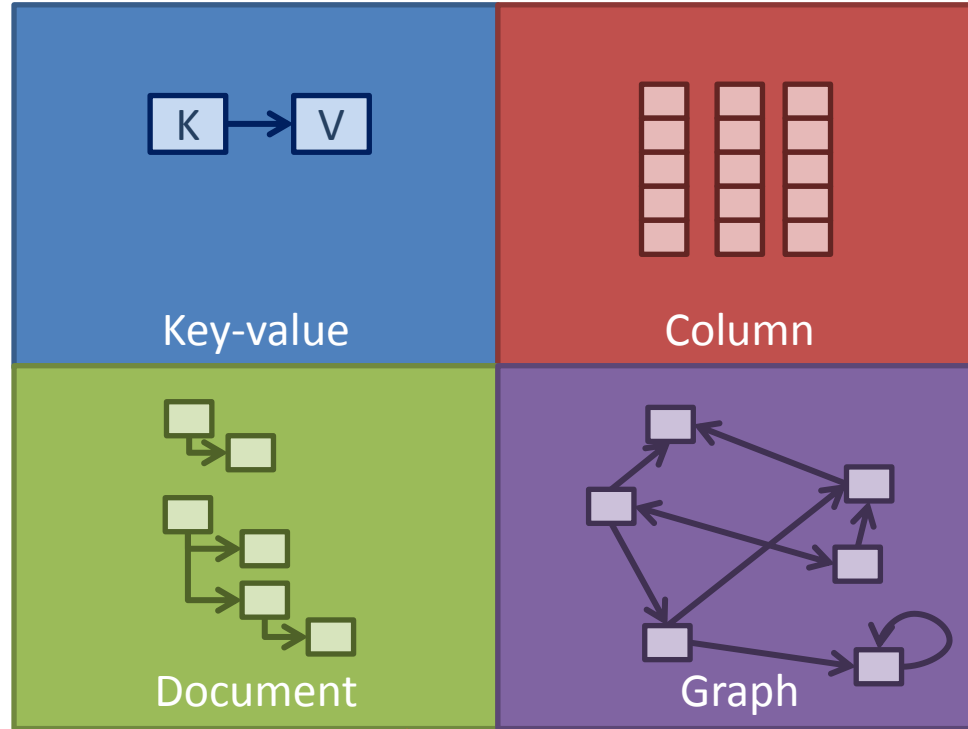
NoSQL solutions



cassandra



ORACLE[®]
NOSQL DATABASE



Google[™]
BigTable

APACHE
HBASE

AllegroGraph

neo4j

InfiniteGraph
Powered by Objectivity

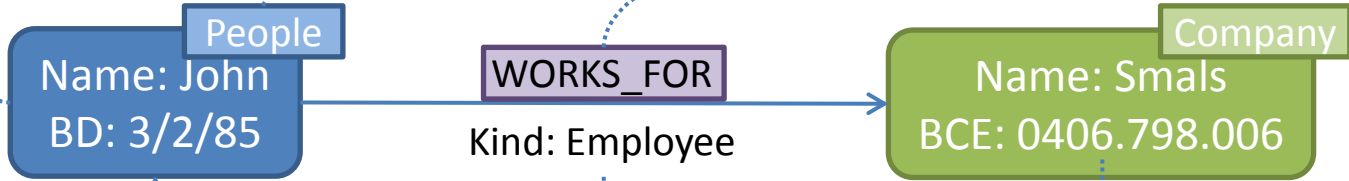
mongoDB[®]

OrientDB[®]

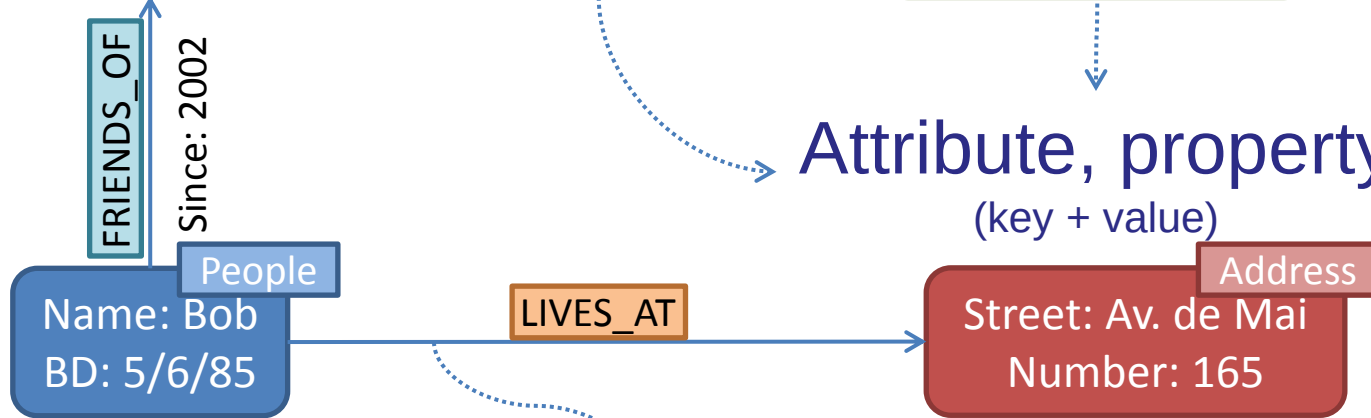
Graph database model

Node, entity

Label, Type



Attribute, property
(key + value)



Relationships

GraphDB: objectives

Objectives of graph databases:

Querying language splitting
apart relationship definition
(creation time) and
« **matching** »

```
SELECT p1.Name
FROM People p1
JOIN Likes l1
  ON l1.Liked_ID = p1.ID
JOIN Likes l2
  ON l1.Liker_ID = l2.Liked_ID
JOIN People p2
  ON l2.Liker_ID = p2.ID
WHERE p2.Name = "Bob"
```

Fast search engine
for relationship
traversals

Cypher (Neo4j)

```
MATCH
  (:People {Name:"Bob"})
-[:Likes*2]->
  (p:People)
RETURN p.Name
```

GraphDB: less code

```
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
FROM (
  SELECT manager.pid AS directReportees, 0 AS count
  FROM person_reportee manager
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  UNION
  SELECT manager.pid AS directReportees, count(manager.directly_manages) AS count
  FROM person_reportee manager
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
  UNION
  SELECT manager.pid AS directReportees, count(reportee.directly_manages) AS count
  FROM person_reportee manager
  JOIN person_reportee reportee
  ON manager.directly_manages = reportee.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
  UNION
  SELECT manager.pid AS directReportees, count(L2Reportees.directly_manages) AS count
  FROM person_reportee manager
  JOIN person_reportee L1Reportees
  ON manager.directly_manages = L1Reportees.pid
  JOIN person_reportee L2Reportees
  ON L1Reportees.directly_manages = L2Reportees.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
) AS T
GROUP BY directReportees)
UNION
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
FROM (
  SELECT manager.directly_manages AS directReportees, 0 AS count
  FROM person_reportee manager
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  UNION
  SELECT reportee.pid AS directReportees, count(reportee.directly_manages) AS count
  FROM person_reportee manager
  JOIN person_reportee reportee
  ON manager.directly_manages = reportee.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
  UNION
  SELECT L2Reportees.pid AS directReportees, count(L2Reportees.directly_manages)
  AS count
  FROM person_reportee manager
  JOIN person_reportee L1Reportees
  ON manager.directly_manages = L1Reportees.pid
  JOIN person_reportee L2Reportees
  ON L1Reportees.directly_manages = L2Reportees.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
) AS T
GROUP BY directReportees)
```

```
SELECT depth1Reportees.pid AS directReportees,
count(depth2Reportees.directly_manages) AS count
FROM person_reportee manager
JOIN person_reportee L1Reportees
ON manager.directly_manages = L1Reportees.pid
JOIN person_reportee L2Reportees
ON L1Reportees.directly_manages = L2Reportees.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
) AS T
GROUP BY directReportees)
UNION
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
FROM(
  SELECT reportee.directly_manages AS directReportees, 0 AS count
  FROM person_reportee manager
  JOIN person_reportee reportee
  ON manager.directly_manages = reportee.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
  UNION
  SELECT L2Reportees.pid AS directReportees, count(L2Reportees.directly_manages)
  AS count
  FROM person_reportee manager
  JOIN person_reportee L1Reportees
  ON manager.directly_manages = L1Reportees.pid
  JOIN person_reportee L2Reportees
  ON L1Reportees.directly_manages = L2Reportees.pid
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
  GROUP BY directReportees
) AS T
GROUP BY directReportees)
```

MATCH

(boss) - [:MANAGES*0..3] -> (sub) ,
(sub) - [:MANAGES*1..3] -> (report)

WHERE boss.name = « John Doe "
RETURN sub.name **AS** Subordinate,
count(report) **AS** Total

GraphDB : less code

"We found Neo4j to be literally thousands of times faster than our prior MySQL solution, with queries that require 10-100 times less code. Today, Neo4j provides eBay with functionality that was previously impossible."



<https://neo4j.com/news/graph-databases-provide-crystal-ball-online-recommendations-wal-mart-ebay/>

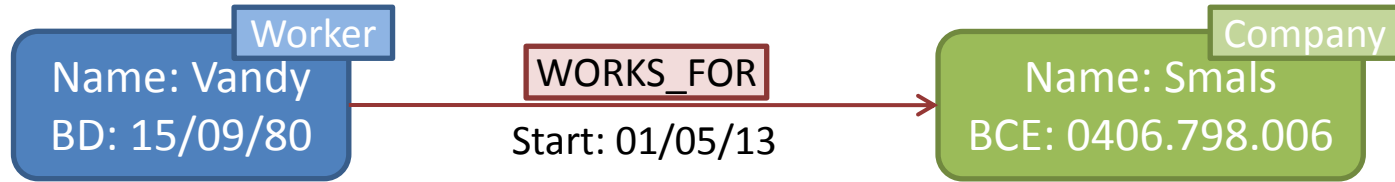
NEO4J

Neo4j

- Current Graph Database leader
- Written in Java, multi-platform
- Community Edition (free, GPL) + Enterprise Edition (\$, AGPL)
- Can be interfaced by many languages
- ACID-compliant
- Queries in Cypher or Gremlin
- Since 2007



Neo4j request



Node of type « **Worker** »

Node of type « **Company** »

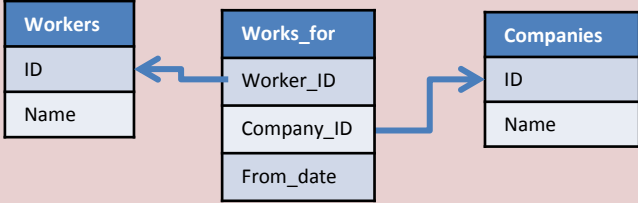
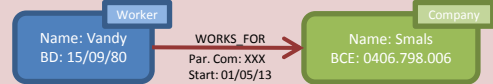
Relationship of type « **WORKS_FOR** »

```
MATCH (w:Worker) -[r:WORKS_FOR]-> (c:Company)
WHERE c.Name = "Smals"
RETURN w.Name, w.BD, r.Start
```

Variante :

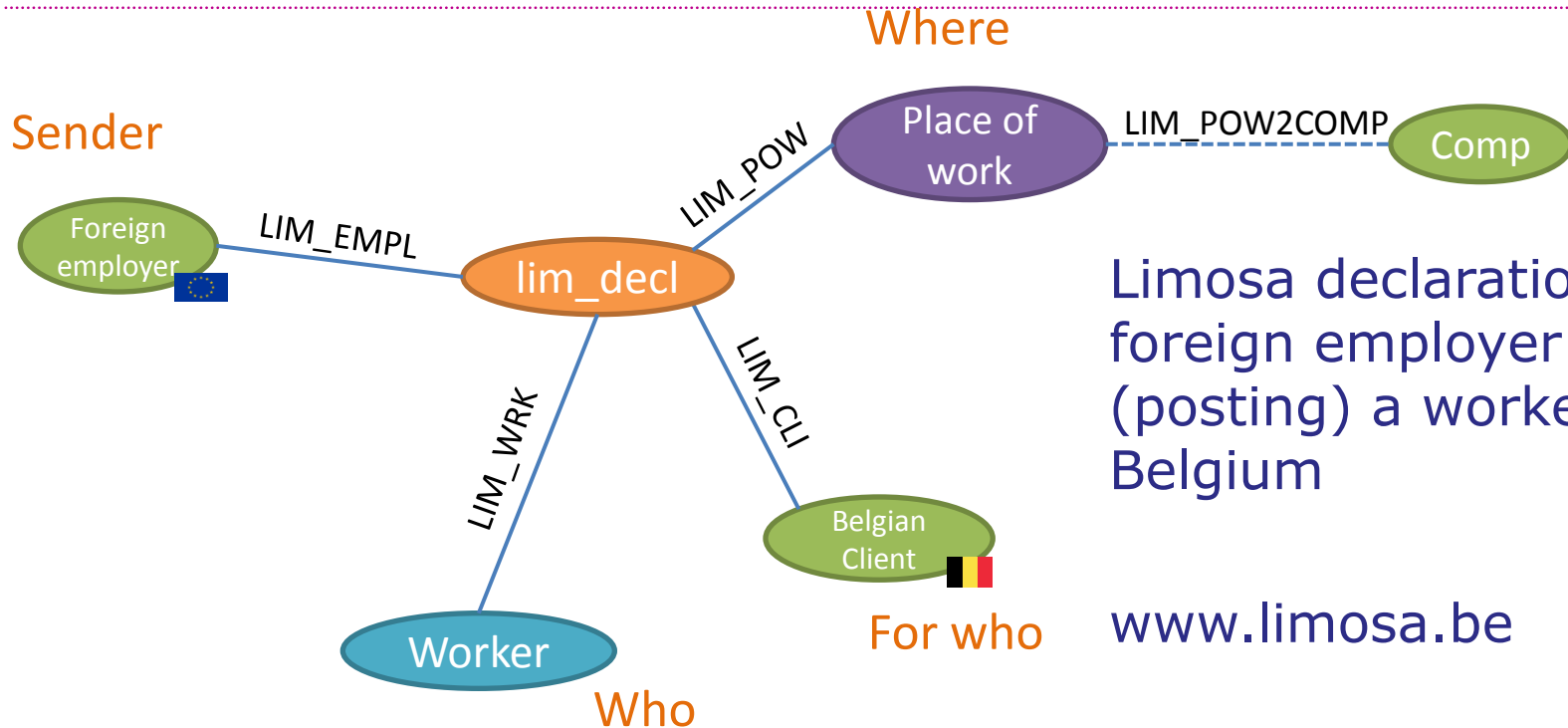
```
MATCH
  (w:Worker)
  -[r:WORKS_FOR]->
  (c:Company {Name:"Smals"})
RETURN ...
```

How it works

	RDBMS	Graph DB
schema		
WHERE ...	Search in Companies (with index): $O(\log([\text{Table size}]))$	Search in nodes (with index) : $O(\log([\text{DB size}]))$
Relation	JOIN : - Search in Works_for : $O(\log([\text{Table size}]))$ - Search in Workers : $O(\log([\text{Table size}]) \times [\text{nb res}])$	Follow pointers in node : $O([\text{nb of employees}])$

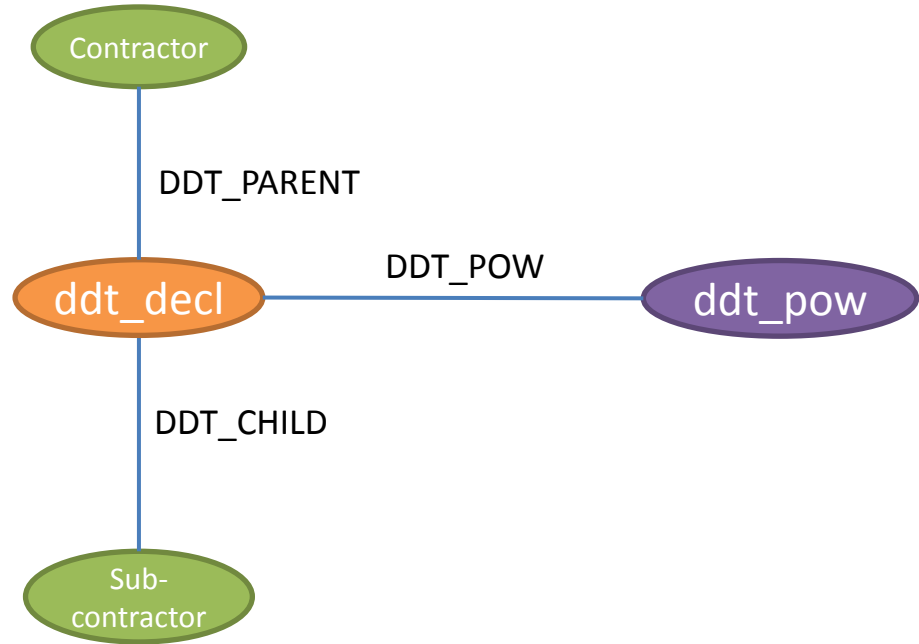
USE CASE: SOCIAL DUMPING

Limosa: detachment declaration



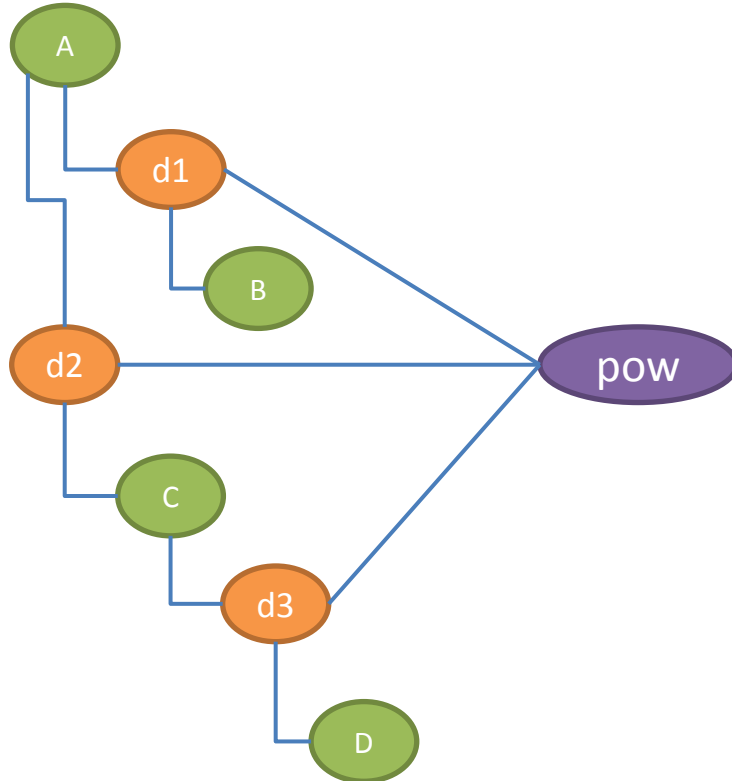
Limosa : *Landenoverschrijdend Informatiesysteem ten behoeve van MigratieOnderzoek bij de Sociale Administratie*
Syst. d'infor. transfrontalier en vue de la recherche en matière de migration auprès de l'adm. sociale

DDT: sub-contractors



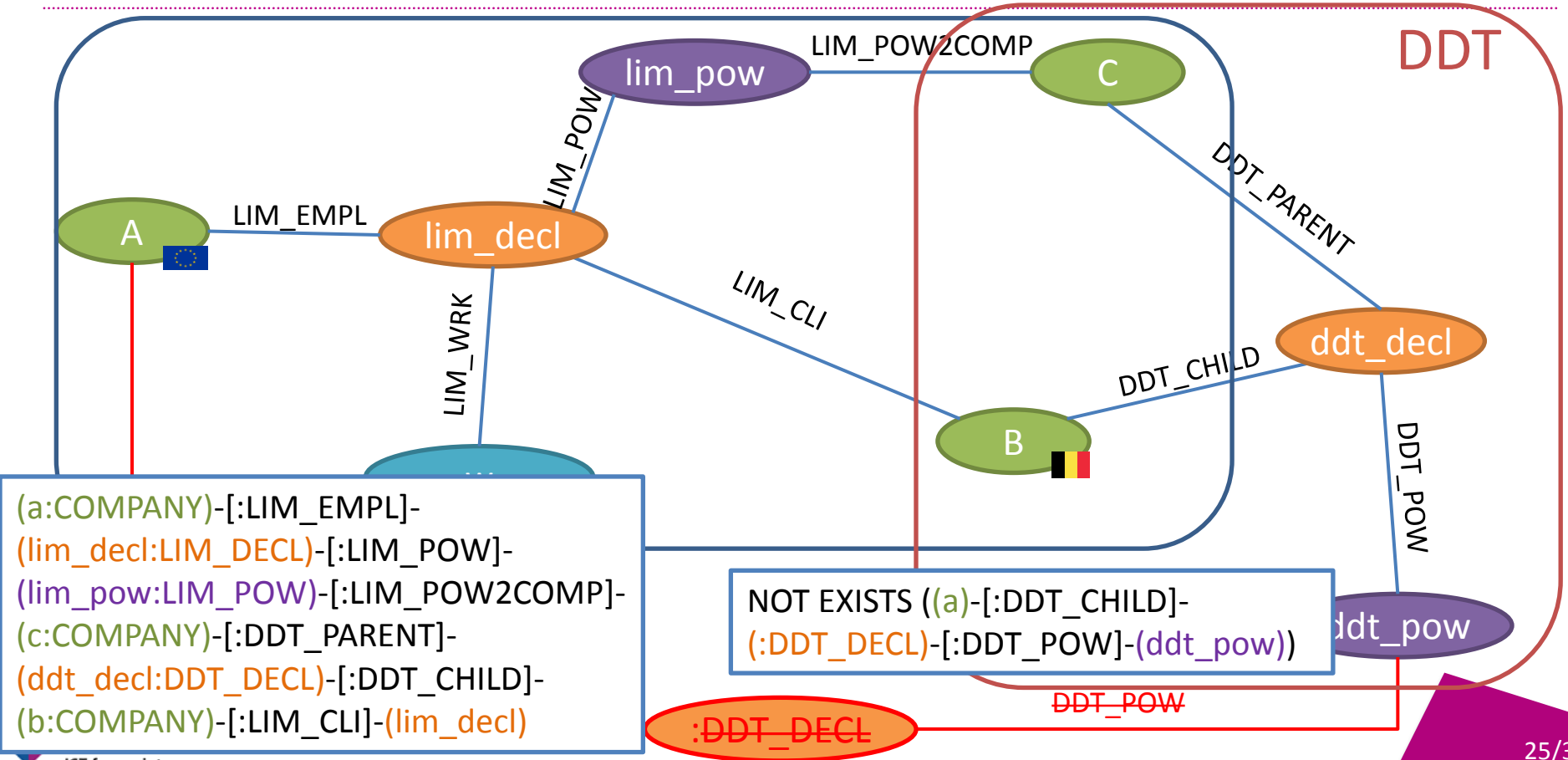
DDT: *D*éclaration *D*e *T*ravaux
Aangifte van werken

DDT: sub-contractors



On a specific worksite (PoW), we have a chain/tree of sub-contractors

Full schema



Full query

MATCH

```
(a:COMPANY)-[:LIM_EMPL]-(lim_decl:LIM_DECL)-[:LIM_POW]-  
(lim_pow:LIM_POW)-[:LIM_POW2COMP]-(c:COMPANY)-[:DDT_PARENT]-  
(ddt_decl:DDT_DECL)-[:DDT_CHILD]-(b:COMPANY)-[:LIM_CLI]-(lim_decl),  
(ddt_decl)-[:DDT_POW]-(ddt_pow:DDT_POW)
```

WHERE

```
NOT EXISTS ((a)-[:DDT_CHILD]-(ddt_decl:DDT_DECL)-[:DDT_POW]-(ddt_pow)) ...
```

RETURN

```
DISTINCT a, b, c,  
COUNT(DISTINCT lim_decl),  
COUNT(DISTINCT ddt_decl)
```

```
ORDER BY COUNT(DISTINCT lim_decl) DESC
```

Implementation

- We copied in Neo4j 17M nodes, 45M relationships :
 - 800.000 companies, 1.650.000 (foreign) workers
 - 13 millions declarations
 - 185.000 duplication groups
 - 250.000 streets
 - ...
- 9 node labels, 15 relationship types
- All companies triplets matching 1st pattern: ~10 seconds (325 results)

In SQL/RDBMS ?

- In SQL/RDBMS,
 - Any node and any relationship is a record : ~65 M records
 - Any type/label corresponds to a table : 24 tables
- Any (n:Type1)-[:REL1]-(m:Type2) corresponds to :

FROM Type1

JOIN REL1 ON Type1.key = REL1.type1_key

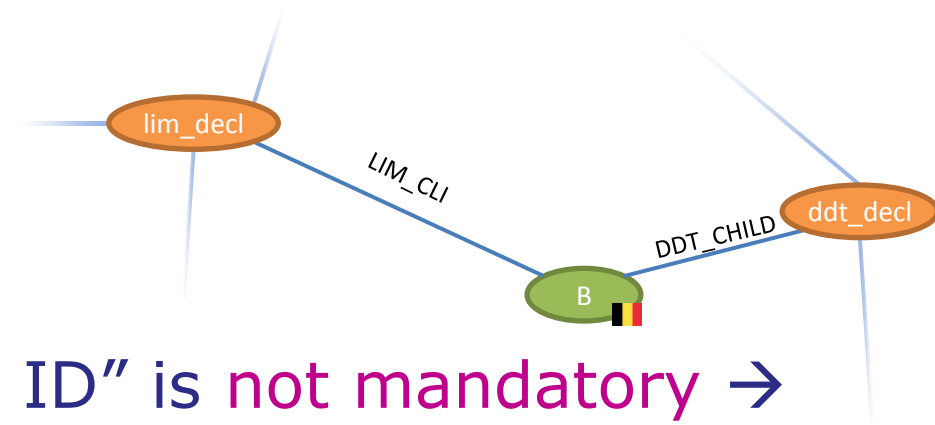
JOIN Type2 ON Type2.key = REL1.type2_key

→ Almost 20 joins!

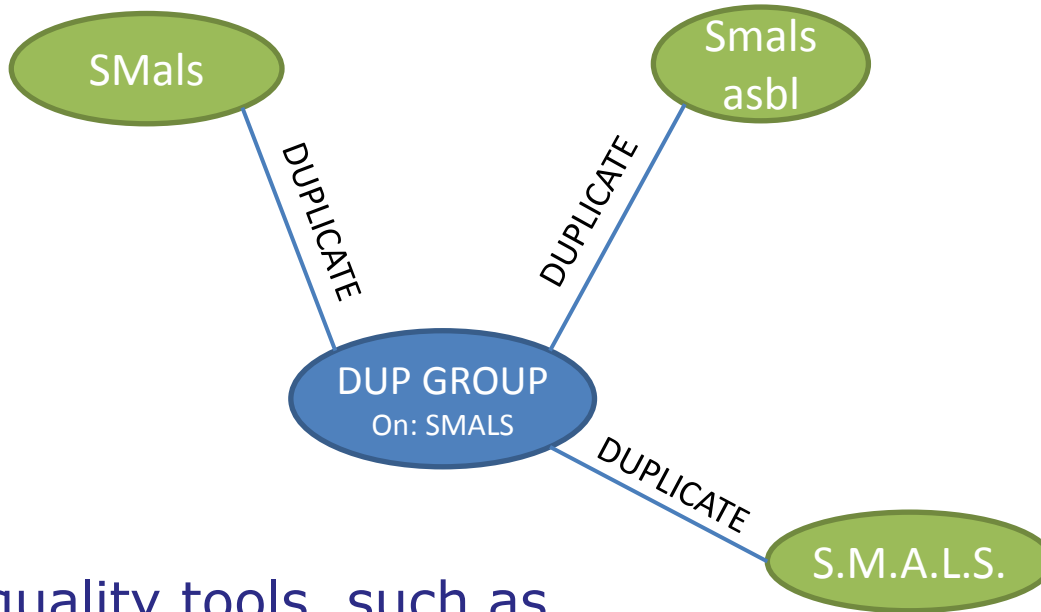
- → Query far more complex to write
- → Much longer to run

Data quality problem

- This only works if « B » has been declared with the **same ID** in both systems !
- Sometimes, the “official ID” is **not mandatory** → name + address is OK
- We have to deal with this!

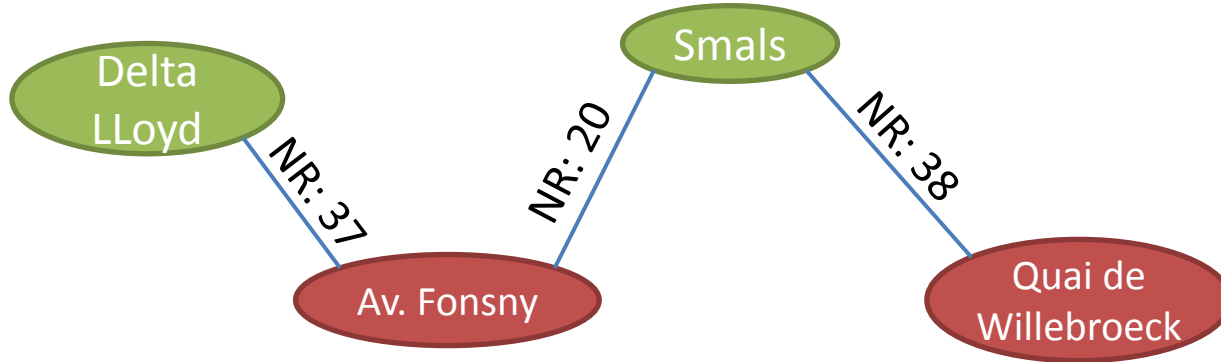


Duplicates



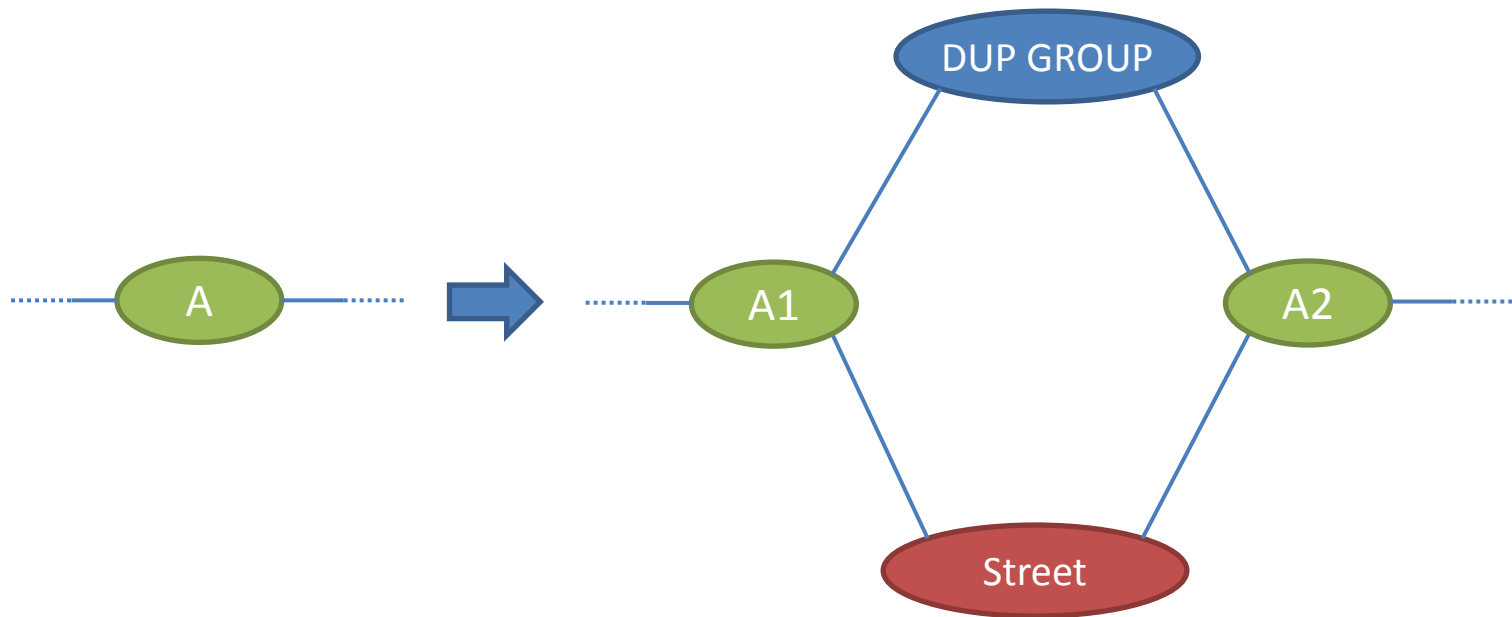
Thanks to data quality tools, such as Trillium (IntoDQ), OpenRefine...

Addresses

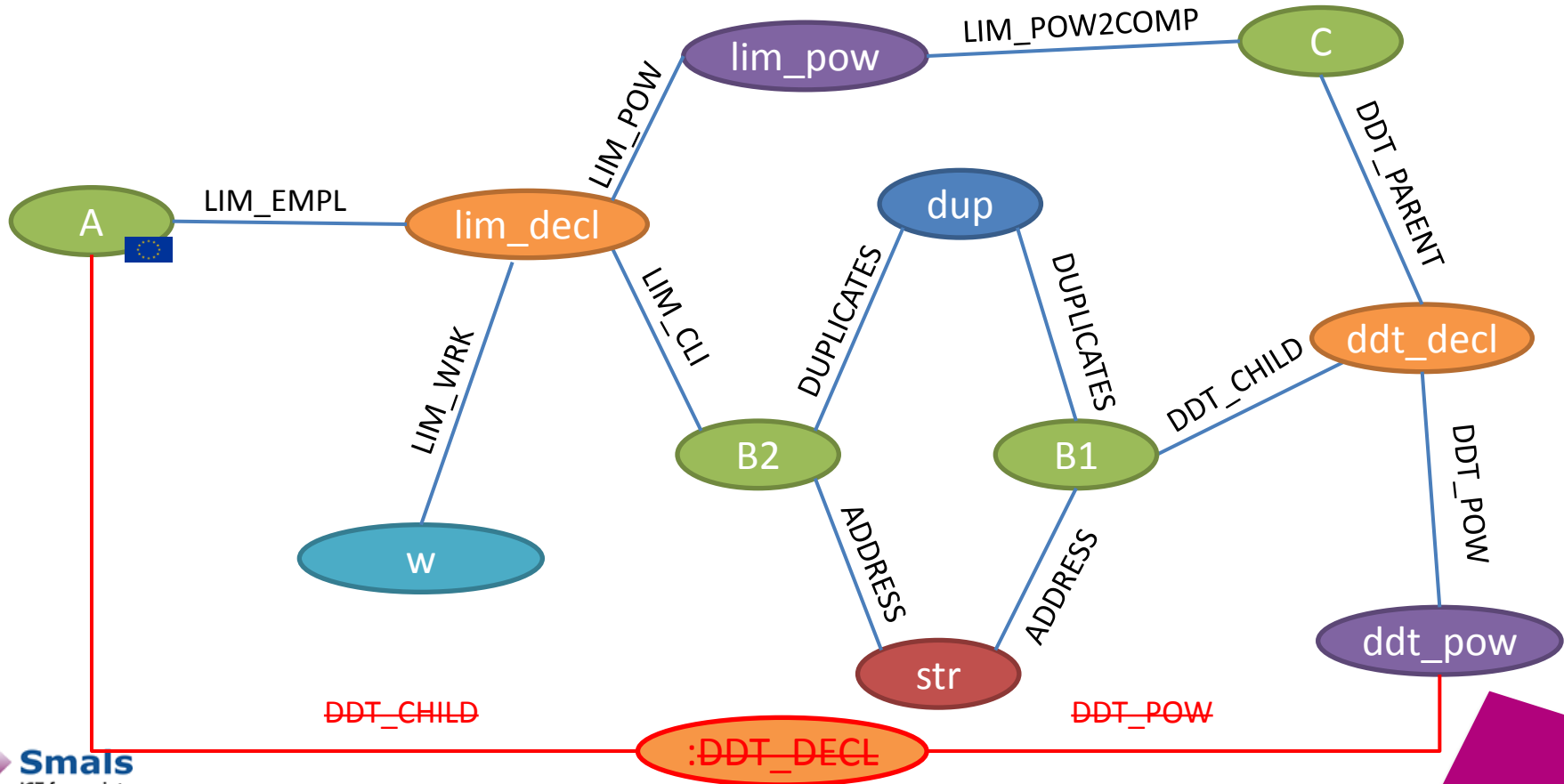


A standardization is needed! Thanks to data quality tools

Duplicates combinations



Schema with B duplicated



Conclusions

- Graph DBs are the « ideal partner » for many fraud detection problems
- They outperform RDBMS on many aspects (query expressiveness, execution speed...)
- Data quality must be taken into consideration
- Not only for fraud: Infrastructure management, recommendation systems, intelligence...

References : www.smalsresearch.be

- Blogs (in French):
 - <https://www.smalsresearch.be/un-fraudeur-ne-fraude-jamais-seul/>
 - <https://www.smalsresearch.be/un-fraudeur-ne-fraude-jamais-seul-partie-2/>
 - <https://www.smalsresearch.be/bases-de-donnees-relationnelles-adequates-pour-des-relations/>
 - <https://www.smalsresearch.be/graph-db-vs-rdbms/>
- Other:
 - NoSQL - Hype ou Innovation, Grégory Ogonowski,
<https://www.smalsresearch.be/publications/document/?docid=59>

References

- **Graph Databases**, *Robinson & all*, O'Reilly 2015
- **Fraud Detection: Discovering Connections with Graph Databases**, *Neo4j*, Sodawksi & Rathle
- **Fraud Analytics Using Descriptive, Predictive, and Social Network Techniques: A Guide to Data Science for Fraud Detection**, *Bart Baesens, Veronique Van Vlasselaer, Wouter Verbeke*, Willey, 2015



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